# Discrete Mathematics and Its Applications 2 (CS147) 

Lecture 14: Chebyshev's inequality and application

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## Recall Markov inequality...

## Statement

Given a non-negative random variable $X$, if its expectation exists, then

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\operatorname{Pr}(X \geq t) \leq \frac{\mathbb{E} X}{t}
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Theorem (Relationship between expectation and tail)
Let $X$ be a non-negative (discrete) random variable taking values in $\{0,1,2, \cdots\}$, if its expectation exists, then

$$
\mathbb{E}(X)=\sum_{i=0}^{\infty} \operatorname{Pr}(X>i) .
$$

## Example: expectation of Geometric distribution (proof by tail)

$X \sim \mathrm{Geo}(p)$ with the PMF

$$
\operatorname{Pr}(X=k)=(1-p)^{k-1} p \quad \forall k \geq 1
$$

## Statement

The expected value of a Geometric random variable is $\mathbb{E}(X)=1 / p$.

## Proof.

Using the integral identity and $q:=1-p$, we have

$$
\begin{aligned}
\mathbb{E}(X) & =\sum_{i=0}^{\infty} \operatorname{Pr}(X>i)=\sum_{i=1}^{\infty} \operatorname{Pr}(X \geq i)=\sum_{i=1}^{\infty} \sum_{k=i}^{\infty}(1-p)^{k-1} p:=p \sum_{i=1}^{\infty} \sum_{k=i}^{\infty} q^{k-1} \\
& =p \sum_{i=1}^{\infty} \frac{q^{i-1}}{1-q}=\sum_{i=1}^{\infty} q^{i-1}=\frac{1}{1-q}=\frac{1}{p}
\end{aligned}
$$

## Recall Variance...

## Definition

The variance of a random variable $X$ is defined as

$$
\mathbb{V}(X)=\mathbb{E}(X-\mathbb{E} X)^{2}=\mathbb{E} X^{2}-[\mathbb{E} X]^{2} .
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## Property

- $\mathbb{V}(a X)=a^{2} \mathbb{V}(X)$ for a constant $a$.
- If $X, Y$ are independent, we have $\mathbb{V}(a X+b Y)=a^{2} \mathbb{V}(X)+b^{2} \mathbb{V}(Y)$.


## Example: Variance of Geometric distribution

$X \sim \operatorname{Geo}(p)$ with the probability mass function

$$
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The variance of a Geometric random variable is $\mathbb{V}(X)=\frac{1-p}{p^{2}}$.

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## Statement

The variance of a Geometric random variable is $\mathbb{V}(X)=\frac{1-p}{p^{2}}$.

## Proof.

We know that $\mathbb{E}(X)=1 / p$ and $\mathbb{V}[X]=\mathbb{E} X^{2}-(\mathbb{E} X)^{2}$, then we only need to know

$$
\begin{aligned}
\mathbb{E}\left(X^{2}\right) & =\sum_{k=1}^{\infty} k^{2}(1-p)^{k-1} p=p \sum_{k=1}^{\infty} k^{2} q^{k-1}=p \sum_{k=1}^{\infty}\left(k q^{k}\right)^{\prime} \quad \text { taking } q:=1-p \\
& =p\left(\sum_{k=1}^{\infty} k q^{k}\right)^{\prime}=p\left(\frac{q}{(1-q)^{2}}\right)^{\prime}=\frac{2-p}{p^{2}}
\end{aligned}
$$

[ $S:=\sum_{k=1}^{\infty} k q^{k}$, using $S-q S=\ldots$ ]

## Chebyshev's inequality

Theorem (Chebyshev's inequality)
For a random variable $X$ with its expectation $\mu$ and variance $\sigma^{2}$, then

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\operatorname{Pr}[|X-\mu| \geq t] \leq \frac{\sigma^{2}}{t^{2}}
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## Proof.

$$
\operatorname{Pr}[|X-\mu| \geq t]=\operatorname{Pr}\left[|X-\mu|^{2} \geq t^{2}\right] \leq \frac{\mathbb{E}\left[|X-\mu|^{2}\right]}{t^{2}}=\frac{\sigma^{2}}{t^{2}} .
$$

More information, better result

- Markov inequality: only use $\mu$, convergence rate: $\mathcal{O}(1 / t)$
- Chebyshev's inequality: use $\mu, \sigma^{2}$, convergence rate: $\mathcal{O}\left(1 / t^{2}\right)$


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More general version: If we introduce a non-decreasing, non-negative function $\phi$, then

$$
\operatorname{Pr}(|X-\mu| \geq t)=\operatorname{Pr}[\phi(|X-\mu|) \geq \phi(t)] \leq \frac{\mathbb{E}[\phi(X)]}{\phi(t)}
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Moment (if exists and finite) by choosing $\phi$ as a polynomial function:

- 1st order moment: $\mathbb{E}[X]$
- 2nd order moment: $\mathbb{E}\left[X^{2}\right], \mathbb{E}\left[|X-\mathbb{E} X|^{2}\right]$
- $t$-th order moment: $\mathbb{E}\left[X^{t}\right], \mathbb{E}\left[|X-\mathbb{E} X|^{t}\right]$


## Application of Chebyshev's inequality to Coupon collector's problem

## Problem (Recall Coupon collector's problem)

We randomly and uniformly sample one object from $\{1,2, \cdots, n\}, T$ is the number of draws before the every $\{1,2, \cdots, n\}$ is seen, we have $\mathbb{E}(T)=n H_{n}$.

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Now we plan to estimate the tail by Chebyshev's inequality.

## Problem (Tail probability)

For coupon collector's problem, what is the probability of the event that the numbers we draw is larger than $N$ ?

$$
\operatorname{Pr}(T \geq N) \leq ?
$$

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- Recall Geometric distribution: $\mathbb{E}(X)=1 / p$ and $\mathbb{V}(X)=\frac{1-p}{p^{2}}$.


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- $T_{i} \sim \operatorname{Geo}\left(p_{i}\right)$ with $p_{i}=\frac{n-i+1}{n}$.
- Recall Geometric distribution: $\mathbb{E}(X)=1 / p$ and $\mathbb{V}(X)=\frac{1-p}{p^{2}}$.
- $\mathbb{E}\left(T_{i}\right)=\frac{n}{n-i+1}$ and $\mathbb{V}\left(T_{i}\right)=\frac{n(i-1)}{(n-i+1)^{2}}$
- $T_{i}$ and $T_{j}$ are independent (each trial is independent)

We estimate the variance of $T$ by the independence of $\left\{T_{i}\right\}_{i=1}^{n}$

## Continue

Solution (To be continued)

$$
\begin{aligned}
\mathbb{V}[T] & =\mathbb{V}\left[\sum_{i=1}^{n} T_{i}\right]=\sum_{i=1}^{n} \mathbb{V}\left(T_{i}\right)=\sum_{i=1}^{n} \frac{n(i-1)}{(n-i+1)^{2}} \\
& =0+\frac{n}{(n-1)^{2}}+\cdots+\frac{n(i-1)}{(n-i+1)^{2}}+\cdots+\frac{n(n-1)}{1^{2}} \\
& =n\left(0+\frac{1}{(n-1)^{2}}+\frac{2}{(n-2)^{2}}+\cdots+\frac{i-1}{(n-i+1)^{2}}+\cdots+\frac{n-1}{1^{2}}\right) \\
& \leq n^{2} \sum_{i=1}^{n} \frac{1}{i^{2}}=\frac{n^{2} \pi^{2}}{6} .
\end{aligned}
$$

## Another way

$$
\begin{aligned}
\mathbb{V}[T] & =n\left(0+\frac{1}{(n-1)^{2}}+\frac{2}{(n-2)^{2}}+\cdots+\frac{i-1}{(n-i+1)^{2}}+\cdots+\frac{n-1}{1^{2}}\right) \\
& \leq n\left(\frac{1}{(n-1)(n-2)}+\frac{2}{(n-2)(n-3)}+\frac{3}{(n-3)(n-4)}+\cdots+\frac{n-2}{2 \times 1}+\frac{n-1}{1^{2}}\right) \\
& =n\left(\left[\frac{1}{n-2}-\frac{1}{n-1}\right]+2\left[\frac{1}{n-3}-\frac{1}{n-2}\right]+3\left[\frac{1}{n-4}-\frac{1}{n-3}\right]+\cdots+(n-2)\left[\frac{1}{1}-\frac{1}{2}\right]+\frac{n-1}{1}\right) \\
& \leq n\left(-\frac{1}{n-1}-\frac{1}{n-2}-\frac{1}{n-3}-\cdots-\frac{1}{2}+(n-2)+(n-1)\right) \\
& =n\left(-\left(H_{n}-1-1 / n\right)+2 n-3\right) .
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\end{aligned}
$$

Similarly, we have the lower bound (using $\frac{1}{n^{2}} \geq \frac{1}{n(n+1)}$ )

$$
\mathbb{V}[T] \geq n\left(-\left(H_{n}-1\right)+n-1\right) .
$$

## Continue

By Chebyshev's inequality, we have

$$
\operatorname{Pr}(|T-\mathbb{E}(T)| \geq t) \leq \frac{\mathbb{V}(X)}{t^{2}} \leq \frac{n^{2} \pi^{2}}{6 t^{2}}
$$

That means,

$$
\operatorname{Pr}(T \geq t+\mathbb{E}(T)) \leq \operatorname{Pr}(|T-\mathbb{E}(T)| \geq t) \leq \frac{n^{2} \pi^{2}}{6 t^{2}}
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Recall $\mathbb{E}(T)=n H_{n}$, we have

$$
\operatorname{Pr}\left(T \geq t+n H_{n}\right) \leq \operatorname{Pr}\left(\left|T-n H_{n}\right| \geq t\right) \leq \frac{n^{2} \pi^{2}}{6 t^{2}}
$$

Taking $t:=(\beta-1) n H_{n}$ with $\beta>1$, we have

$$
\operatorname{Pr}\left(T \geq \beta n H_{n}\right) \leq \frac{\pi^{2}}{6(\beta-1)^{2} H_{n}^{2}} \leq \frac{\pi^{2}}{6(\beta-1)^{2} \log ^{2} n}
$$

## Can we do it better?

$$
\operatorname{Pr}(T \geq N) \leq \text { small }
$$

- strictly speaking, it should be $T \geq N+1 \ldots$
- that means, at least one of $n$ distinct objects has not been selected in the first $N$ round.


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- Let $A_{i}^{N}$ denote when item $i$ is not observed in the first $N$ draws, i.e., $\operatorname{Pr}\left(A_{i}^{N}\right)=\left(1-\frac{1}{n}\right)^{N}$.
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Taking $N:=\beta n \log n$, we have (using $1+x \leq e^{x}$ for any $x \in \mathbb{R}$ )

$$
n\left(1-\frac{1}{n}\right)^{N} \leq n\left(e^{-\frac{1}{n}}\right)^{\beta n \log n}=n e^{-\beta \log n}=n\left(e^{\log n}\right)^{-\beta}=n^{-\beta+1} .
$$

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$$

$\Rightarrow \operatorname{Pr}(T \geq N) \leq n^{-\beta+1}$.

## References I

[0] Tao Luo, Zhi-Qin John Xu, Zheng Ma, and Yaoyu Zhang, Phase diagram for two-layer relu neural networks at infinite-width limit, Journal of Machine Learning Research 22 (2021), no. 71, 1-47.
(Cited on pages 31 and 32 .)

## *Application in ML theory: tail and union bound [LXMZ21]

Lemma 9 (bounds of initial parameters). Given $\delta \in(0,1)$, we have with probability at least $1-\delta$ over the choice of $\boldsymbol{\theta}^{0}$

$$
\begin{equation*}
\max _{k \in[m]}\left\{\left|a_{k}^{0}\right|,\left\|\boldsymbol{w}_{k}^{0}\right\|_{\infty}\right\} \leq \sqrt{2 \log \frac{2 m(d+1)}{\delta}} \tag{45}
\end{equation*}
$$

Proof If $\mathrm{X} \sim N(0,1)$, then $\mathbb{P}(|\mathrm{X}|>\varepsilon) \leq 2 \mathrm{e}^{-\frac{1}{2} \varepsilon^{2}}$ for all $\varepsilon>0$. Since $a_{k}^{0} \sim N(0,1)$, $\left(w_{k}^{0}\right)_{\alpha} \sim N(0,1)$ for $k=1,2, \ldots, m, \alpha=1, \ldots, d$ and they are all independent, by setting

$$
\varepsilon=\sqrt{2 \log \frac{2 m(d+1)}{\delta}}
$$

one can obtain

$$
\begin{aligned}
\mathbb{P}\left(\max _{k \in[m]}\left\{\left|a_{k}^{0}\right|,\left\|\boldsymbol{w}_{k}^{0}\right\|_{\infty}\right\}>\varepsilon\right) & =\mathbb{P}\left(\max _{k \in[m], \alpha \in[d]}\left\{\left|a_{k}^{0}\right|,\left|\left(w_{k}^{0}\right)_{\alpha}\right|\right\}>\varepsilon\right) \\
& =\mathbb{P}\left(\bigcup_{k=1}^{m}\left(\left|a_{k}^{0}\right|>\varepsilon\right) \bigcup\left(\bigcup_{\alpha=1}^{d}\left(\left|\left(w_{k}^{0}\right)_{\alpha}\right|>\varepsilon\right)\right)\right) \\
& \leq \sum_{k=1}^{m} \mathbb{P}\left(\left|a_{k}^{0}\right|>\varepsilon\right)+\sum_{k=1}^{m} \sum_{\alpha=1}^{d} \mathbb{P}\left(\left|\left(w_{k}^{0}\right)_{\alpha}\right|>\varepsilon\right) \\
& \leq 2 m \mathrm{e}^{-\frac{1}{2} \varepsilon^{2}}+2 m d \mathrm{e}^{-\frac{1}{2} \varepsilon^{2}} \\
& =2 m(d+1) \mathrm{e}^{-\frac{1}{2} \varepsilon^{2}} \\
& =\delta
\end{aligned}
$$

## *Application in ML theory: Markov inequality [LXMZ21]

Then

$$
\begin{aligned}
\mathbb{E} \sum_{i, j=1}^{n} & \left|G_{i j}^{[\boldsymbol{w}]}(\boldsymbol{\theta}(t))-G_{i j}^{[\boldsymbol{w}]}(\boldsymbol{\theta}(0))\right| \\
& \leq \sum_{i, j=1}^{n} \frac{\kappa^{2} \kappa^{\prime} d}{m} \sum_{k=1}^{m}\left(4 \max \left\{\frac{1}{\kappa^{\prime 2}}, 1\right\} \xi^{2} \mathbb{E}\left|D_{k, i, j}\right|+6 \max \left\{\frac{1}{\kappa^{\prime 2}}, 1\right\} \xi^{2} p\right) \\
& \leq \sum_{i, j=1}^{n} \frac{\kappa^{2} \kappa^{\prime} d}{m} \sum_{k=1}^{m}\left(4 \max \left\{\frac{1}{\kappa^{\prime 2}}, 1\right\} \xi^{2} 8 d \max \left\{\kappa^{\prime}, 1\right\} \xi p+6 \max \left\{\frac{1}{\kappa^{\prime 2}}, 1\right\} \xi^{2} p\right) \\
& \leq \kappa^{2} \kappa^{\prime} d n^{2}\left(32 d \xi \max \left\{\kappa^{\prime}, \frac{1}{\kappa^{\prime 2}}\right\}+6 \max \left\{\frac{1}{\kappa^{\prime 2}}, 1\right\}\right) \xi^{2} p \\
& \leq 40 \kappa^{2} d^{2} n^{2}\left(2 \log \frac{8 m(d+1)}{\delta}\right)^{3 / 2} \max \left\{\kappa^{\prime 2}, \frac{1}{\kappa^{\prime}}\right\} p .
\end{aligned}
$$

By Markov's inequality, with probability at least $1-\delta / 2$ over the choice of $\boldsymbol{\theta}^{0}$, we have

$$
\begin{aligned}
& \left\|G^{[\boldsymbol{w}]}(\boldsymbol{\theta}(t))-G^{[\boldsymbol{w}]}(\boldsymbol{\theta}(0))\right\|_{\mathrm{F}} \\
& \leq \sum_{i, j=1}^{n}\left|G_{i j}^{[\boldsymbol{w}]}(\boldsymbol{\theta}(t))-G_{i j}^{[\boldsymbol{w}]}(\boldsymbol{\theta}(0))\right| \\
& \leq \max \left\{\kappa^{\prime 2}, \frac{1}{\kappa^{\prime}}\right\} \frac{40 \kappa^{2} d^{2} n^{2}\left(2 \log \frac{8 m(d+1)}{\delta}\right)^{3 / 2} p}{\delta / 2}
\end{aligned}
$$

